

Stochastic Modeling, Algorithms and Analysis for Consensus Seeking over Noisy Networks

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Detroit, April 2008

Contents

- Background: from animal behavior to engineering
- Existing research
- Consensus seeking in uncertain environment
- Stochastic algorithms
- Convergence and performance
- Concluding remarks

Animal Behavior: Birds

• A group of **birds** fly with coordination in speed and direction (**Flocking**)





Fish

 Huge number of fish cooperatively move (Schooling)

-- Important for search for food or for protection from predators



Couzin et.al. *Nature*, 2005

Honeybees

• Honeybees select a new home from several candidate sites spotted by scout bees

 What is the mechanism for reaching consensus? (Visscher, *Nature*, 2003)

-- Important for avoiding population disintegration



From Birds to Bees: from Flocking/Swarming to Consensus

- Each agent has local information about neighboring agents
- and there is a key group objective (e.g., achieve accurate alignment during motion, or agree on a nest site, etc.)

Such coordination amounts to a form of consensus

Math theory for interpretation?



Applications in Technology

- Examples: a group of autonomous vehicles, or robot teams (formation control)
- In such distributed multi-agent control systems coordination is critical for safety & the performance of tasks (below: simple robots)





Formation of Platoon of Vehicles

- Equalize velocity of different vehicles
- Maintain spacing
- Increase highway capacity and improve safety



The Consensus Issue

- For multi-agent coordination, it is usually important to maintain shared information between agents
- This leads to the key issue of "Agreeing-on-something". This agreement may
 - (1) be the **objective of operation**
 - (2) or a **condition** for proceeding to further operation

Hence, in this context, we study consensus problems.

What Is Consensus?

- By consensus seeking, we mean a mechanism whereby the agents adjust their individual values of an underlying quantity (e.g., a key state value – angle, velocity, etc.) so as to converge to a common value
- In general, convergence is a primary objective
- The actually reached limit may be of secondary importance

(small fish schooling)



Background: Models with Exact State Info

- Most existing research on consensus problems assumes exact state information exchange
- Maintaining certain connectivity (which can be relaxed to different forms) is crucial for achieving consensus



• The most important analytical tools come from the theory of stochastic matrices

Background: Models with Noisy or Inaccurate Measurements

- In a distributed network, it may be impractical to have exact state exchange, for example, due to --- receiver noise
 - --- quantization, etc. etc.
- Consensus models with additive noises have attracted the interest of many authors
 - --- (Ren, Beard and Kingston, ACC'05)
 - --- (Xiao, Boyd, and Kim, 2007)

--- (Huang and Manton, ACC'07, CDC'07, ACC'08, Preprint'06, Preprint'08)

--- More recent works by various authors ...

Related stochastic models for consensus

--- (Tsitsiklis, Bertsekas, and Athens, IEEE TAC'86) stochastic gradient based algorithms for distributed function optimization



Definitions

 Definition 1 (weak consensus) The agents are said to reach weak consensus if

$$\lim_{t \to \infty} E|x_t^i - x_t^j|^2 = 0, \quad \forall i, j.$$

• Definition 2 (mean square consensus) The agents are said to reach m.s. consensus if $E|x_t^i|^2 < \infty$, $\forall i \in \mathcal{N}, t$ and there exists x^* such that

$$\lim_{t \to \infty} E|x_t^i - x^*|^2 = 0, \ \forall i \in \mathcal{N}$$

• Definition 3 (strong consensus) The agents are said to reach strong consensus if there exists x^* such that

 $x_t^i
ightarrow x^*$ with probability one for all i.

Graph Modeling of Networked Agents

- Consider directed graphs (i.e., digraphs) $G = (\mathcal{N}, \mathcal{E})$
- Each agent is denoted by a node
- In a digraph, arrow indicates neighboring relationship & infor. flow (Example -- right top, agent 1 is a neighbor of agent 2)
- In undirected graph (special case), information is bidirectional



Network Topology Modeling

 For our further analysis: we assume---

The digraph contains a spanning tree (special case: connected undirected graphs)

- Implication: information may propagate across the network from one or more points
- In a deterministic model with fixed topology, Ren et. al. (2005) proved existence of a spanning tree is the weakest connectivity condition for consensus



The Measurement Model

- Each agent knows its own state x_t^i exactly,
- and it has noisy observation y_t^{ik} of its neighbors' states, i.e.,

$$y_t^{ik} = x_t^k + w_t^{ik}, \qquad t \in Z^+, \quad k \in \mathcal{N}_i.$$

where w_t^{ik} is additive measurement noise.



If Fixed Coefficients Are Used in Averaging: What Happens?

$$\begin{cases} x_{t+1}^1 = \frac{1}{2}(x_t^1 + y_t^{12}) \\ x_{t+1}^2 = \frac{1}{3}(x_t^2 + y_t^{21} + y_t^{23}) \\ x_{t+1}^3 = \frac{1}{2}(x_t^3 + y_t^{32}) \end{cases}$$

This algorithm is essentially a noisy variant of equal-neighbor based algorithms (see related algorithms: Vicsek et. al. PRL'95 Jadbabaie, Lin, Morse'03, etc.)

Measurement noise causes divergence.





Stochastic Approximation

• We use the averaging rule (convex combination):

$$\begin{aligned} x_{t+1}^i &= (1 - a_t b_{ii}) x_t^i + a_t \sum_{k \in \mathcal{N}_i} b_{ik} y_t^{ik}, \quad t \ge 0 \\ b_{ik} &> 0 \text{ if and only if } k \in \mathcal{N}_i \\ b_{ii} &= \sum_{k \in \mathcal{N}_i} b_{ik} \end{aligned}$$

• The state of a node remains the same if it has no neighbors. (This happens in leader following)

Stochastic Approximation

• The algorithm in vector form:

$$x_{t+1} = x_t + a_t B x_t + a_t \tilde{w}_t$$

where B has zero row sum.

$$B = \begin{bmatrix} -b_{11} & b_{12} & \cdots & b_{1n} \\ b_{21} & -b_{22} & \cdots & b_{2n} \\ \vdots & & & \vdots \\ b_{n1} & b_{n2} & \cdots & -b_{nn} \end{bmatrix}$$

where

$$b_{ij} = 0$$
 if $j \notin \mathcal{N}_i \cup \{i\}$.

• *B* is unstable and may be viewed as the generator of a continuous time Markov chain.

Main Assumptions

- (A1) The measurement noises are independent random variables with zero mean, and independent of initial states.
 The noise and initial states have bounded second order moment.
- (A2) The digraph contains a spanning tree.
- (A3) The positive step size sequence $\{a_t, t \ge 0\}$ satisfies:

$$\sum_{i=0}^{\infty} a_i^2 < \infty, \quad \sum_{i=0}^{\infty} a_i = \infty$$

Remark: The independence noise sequence assumption may be relaxed (for instance, a sequence of martingale differences for noise vectors)

Illustration with a Two-agent Model

• First, under (A1)-(A2) for noise and step size, it is relatively easy to show (a.s. and m.s.) convergence of the mid-point

$$z_t = \frac{1}{2}(x_t^1 + x_t^2) \to z^*$$

- Next it suffices to show (a.s. and m.s.) convergence of the state gap $\xi_t = x_t^1 x_t^2$
- We have $\xi_{t+1} = (1-2a_t)\xi_t + a_t v_t, \quad t \geq 0$

where $v_t = w_t^{12} - w_t^{21}$

The Diagram for State Gap



 Key idea: show benefits of reducing noise outweigh the disadvantage of reducing stability

State Gap as Noise Summation

• Denote
$$\bar{a}_t = 2a_t$$
 and
 $\Pi_{l,k} = (1 - \bar{a}_l)(1 - \bar{a}_{l-1}) \cdots (1 - \bar{a}_{k+1})a_k$
for $l > k \ge T_1$. We set $\Pi_{k,k} = a_k$.
• The state gap satisfies
 $\xi_{t+1} = (1 - \bar{a}_t)(1 - \bar{a}_{t-1}) \cdots (1 - \bar{a}_{T_1})\xi_{T_1}$
 $+\Pi_{t,T_1}v_{T_1}$
 \vdots
 $+\Pi_{t,t-1}v_{t-1}$
 $+\Pi_{t,t}v_t$
• To prove vanishing gap: Show $\Pi_{t,k}$ or related terms sufficiently small

Convergence Analysis

- Mean square convergence
- Sample path convergence

How to Prove M.S. Convergence?

- Use stochastic Lyapunov analysis to show all individual states attract to each other in mean square
- Next, show the individual states actually go to the same limit.

The Lyapunov Function

- Let S^{n×n} be the set of symmetric matrices and denote
 D = {D ∈ S^{n×n} : D ≥ 0, Null(D) = span{1_n}}
- Lemma. Under (A2) and given $D \in \mathcal{P}$ the

Degenerate Lyapunov Eqn: $QB + B^TQ = -D$ has a unique solution

$$Q \in \mathcal{D}.$$

• The idea is to show the energy function $V(t) = Ex_t^T Qx_t$ will decay to zero.

Energy Decay and Weak Consensus

Theorem (weak consensus). Under (A1)-(A3),
(i) There exist c₁ > 0, c₂ > 0, and a large T_c > 0 such that

$$V(t+1) \le (1 - a_t c_1 + a_t^2 c_2)V(t) + O(a_t^2)$$

(ii) Consequently $\lim_{t \to \infty} V(t) = 0$, which implies

$$\lim_{t \to \infty} E|x_t^i - x_t^k|^2 = 0, \forall i, k.$$

Stay in $span\{1_n\}$! Remark: Here it is not clear yet whether they will

converge to the same limit. (so, need an extra step!)

Mean Square Consensus

• Lemma. There is a unique probability measure π such that $\pi^T B = 0$. Further $\pi^T x_{t+1} = \pi^T x_t + a_t \pi^T \tilde{w}_t$

and $\pi^T x_t$ converges in m.s.

This Lemma combined with $\lim_{t\to\infty} E|x_t^i - x_t^k|^2 = 0, \forall i, k.$ Theorem. (A1)-(A3) ensures Mean Square CONSENSUS (Huang and Manton, ACC'07,08)

Simulations



Further Extension to Leader Following

• For leader following, the stochastic Lyapunov analysis is applicable to establish mean square convergence of all other agents' states to that of the leader (i.e., 4 below).



Left: use direct averaging
 2-0

Right: use stochastic approx.

Sample Path Behavior

- What is the group behavior along sample paths?
- In fact, this can be characterized by sample path convergence

Sample Path Convergence

- Theorem 1. Under (A1)-(A3), the Stochastic Approx. (SA) algorithm ensures strong consensus (i.e. sample path convergence).
- Remark: for strong consensus, the second order moment condition for the noise may be relaxed

Sample Path Analysis via Change of Coordinates

• By choosing a suitable change of coordinates $z_t = [z_t^1, z_t^{(n-1)}]^T = \Phi^{-1}x_t$, the consensus algorithm may be decomposed into the form (Huang & Manton, ACC'08)

$$\begin{cases} z_{t+1}^1 = z_t^1 + a_t v_t^1 \\ z_{t+1}^{(n-1)} = (I + a_t \tilde{B}_{n-1}) z_t^{(n-1)} + a_t v_t^{(n-1)} \end{cases}$$

All eigenvalues of \tilde{B}_{n-1} have negative real parts Thus, we only need to deal with a random walk and a stable linear SA model Alternative Proving Tool: Double Array Analysis

Theorem (Teicher,1985). Let {w, w_t, t ≥} be i.i.d. r.v.'s with zero mean and variance Q and {a_{ki}, 1 ≤ i ≤ l_k ↑ ∞, k ≥ 1}a double array of constants. Assume

i)
$$\max_{1 \le i \le l_k} |a_{ki}| h_i = O(1/\log k)$$
, where $0 < h_i \uparrow h_i = O(i^{1/\delta})$ for some $\delta \in [1, 2]$

(ii)

$$\sum_{i=1}^{\infty} P\{|w| > h_i\} < \infty$$
(iii)

$$h_i/i \downarrow, \text{ and } \sum_{i=1}^{l_k} |a_{ki}|^2 h_i^{2-\delta} = o(1/\log k),$$

$$\sum_{i=1}^{l_k} |a_{ki}|^2 h_i^{2-\delta} = O(1/\log l_k)$$

Then

$$\lim_{k \to \infty} \sum_{i=1}^{l_k} a_{ki} w_i = 0, \qquad a.s.$$

Performance?

Performance Analysis

Due to consensus, denote the limit of the state vector by

$$x_{\infty} = [x_{\infty}^1, \cdots, x_{\infty}^n]^T = x_{\infty}^1 \mathbf{1}_n$$

• Convergence rate --- Roughly, how small is the error term $x_t - x_\infty$ when t is large?

Performance (w/ spanning tree model)

- How fast to reach consensus?(charctrzd by asy. normal.)
- Take step size $a_t = a/t$. Denote $x_t = [x_t^1, \cdots, x_t^n]^T$
- Then under quite standard conditions for noise & coeffic. matrix, we show consensus and furthermore:

$$x_t = x_\infty^1 \mathbf{1}_n + x_t^{e,a} + x_t^{e,b}$$

where $x_t^{e,a}$ depends on future noises & $x_t^{e,b}$ is linear in \mathcal{X}_t $\sqrt{t}x_t^{e,a} \xrightarrow{d} N(0,Q_a), \qquad \sqrt{t}x_t^{e,b} \xrightarrow{d} N(0,Q_b)$

• (H.&M., ACC'08; H. CDC'08 sub) so error decays by rate $\frac{1}{\sqrt{t}}$

Illustration of Asymptotic Normality

Additional Uncertainty Factors

- Random communication link failures
- Quantization effects

Random Link Failures

- The stochastic algorithm may still be applied for the randomly varying topology.
- In this case, the coefficient matrix in the consensus algorithm is given as a sequence of random matrices B_t with mean B

Random Link Failures (ctn)

• The consensus algorithm

$$x_{t+1} = x_t + a_t B_t x_t + a_t \text{``noise''}$$
$$= x_t + a_t \overline{B} x_t + a_t (B_t - \overline{B}) x_t + a_t \text{``noise''}$$

- This algorithm may be viewed as the standard one (with fixed topology) subject to unbiased perturbations.
- In particular, for i.i.d. link failures with additive measurement noise, a perturbed Lyapunov analysis may be applied to establish convergence (Huang and Manton, ACC'08, and Preprint (submitted to IEEE, June'07))

Quantized Data---How to Achieve Convergence?

Probabilistic Quantization (PQ)

- Suppose the state x_t^i is between two quantization levels $r_k < r_{k+1}$
- The idea of PQ is to produce a randomized output Q_i(t) at the quantizer such that it takes the lower and upper level with probability

$$p_k = (r_{k+1} - x_t^i)/(r_{k+1} - r_k), \ p_{k+1} = 1 - p_k$$

respectively

Probabilistic Quantization (PQ)

- This approach has been successfully applied for:
- sensor network signal processing (Xiao, Cui, Luo, and Goldsmith, 2006), and
- **CONSENSUS MODELS** (Aysal, Coates and Rabbat, 2007)

PQ Combined with SA

- In PQ, we may view and quantization error as an additive uncorrelated noise.
- In the consensus algorithm, a decreasing step size may be further used to damp out the noise. Convergence results may be proved. (Huang, Dey, Nair, and Manton, CDC'08 submitted)
- Left: deterministic quantization; Right: PQ

Concluding Remarks

• Stochastic consensus and convergence

- The key is a decreasing step size for cautious learning
- Stochastic Lyapunov analysis is useful
- Many application opportunities in sensor network signal processing (see, e.g. S. Boyd, J. Hespanha) – networked estim. Prob., sensornet time synchronization, sensornet localization etc. etc. etc.

Many practical modeling scenarios -- high order (inertia) models and asynchronous algorithms, approximate average consensus, etc. etc.